

Pointwise Approximation of a Stochastic Heat Equation with Additive Space-Time White Noise

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Stochastic Heat Equation

$$\begin{aligned}dX(t) &= \Delta X(t) dt + B(t) dW(t), \quad t \in [0, T], \\X(0) &= x_0,\end{aligned}$$

with

- ▶ $x_0 \in H = L_2((0, 1))$,
- ▶ Dirichlet boundary conditions,
- ▶ multiplication operator B , i.e.
 $(B(t)h)(u) = g(t, u) \cdot h(u)$
for $u \in [0, 1]$, $h \in H$ and $g \in C^{(1,1)}([0, T] \times [0, 1])$,
- ▶ cylindrical Brownian motion $W = (\langle W(t), h \rangle)_{t \geq 0, h \in H}$.

Computational Problem

Task

Approximate $X(T)$ based on evaluations of finitely many scalar Brownian motions β_i 's at a finite number of points, where

$$\beta_i(\cdot) = \langle W(\cdot), h_i \rangle$$

with $h_i(u) = \sqrt{2} \sin(i\pi u)$ and $i \in \mathbb{N}$.

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Error and cost of approximation $\hat{X}(T)$

$$\begin{aligned} e(\hat{X}(T)) &= \left(\mathbb{E} \|X(T) - \hat{X}(T)\|_H^2 \right)^{1/2}, \\ \text{cost}(\hat{X}(T)) &= \text{total number of evaluations of the } \beta_i \text{'s.} \end{aligned}$$

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Goal

Approximation with optimal relation between error and cost.

Classes of Approximations and Minimal Errors

Class $\mathfrak{X}_{\text{uni}}$ of algorithms with **uniform** time discretization:

- ▶ choose finite set $\mathcal{I} \subset \mathbb{N}$: only evaluate β_i with $i \in \mathcal{I}$.
- ▶ choose $n \in \mathbb{N}$: number of evaluations for β_i with $i \in \mathcal{I}$.
- ▶ choose nodes $t_k = \frac{k}{n}T$, $k = 1, \dots, n$, for β_i with $i \in \mathcal{I}$.
- ▶ choose $\phi : \mathbb{R}^N \rightarrow H$ measurable with $N = |\mathcal{I}| \cdot n$.

Approximation:

$$\widehat{X}(T) = \phi(\beta_{i_1}(t_1), \dots, \beta_{i_1}(t_n), \dots, \beta_{i_\ell}(t_1), \dots, \beta_{i_\ell}(t_n))$$

for $\mathcal{I} = \{i_1, \dots, i_\ell\}$.

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Clearly

$$\text{cost}(\widehat{X}(T)) = N.$$

N -th minimal error

$$e_{\text{uni}}(N) = \inf \left\{ e(\widehat{X}(T)) ; \text{cost}(\widehat{X}(T)) \leq N \text{ and } \widehat{X}(T) \in \mathfrak{X}_{\text{uni}} \right\}.$$

Classes of Approximations and Minimal Errors

Class $\mathfrak{X}_{\text{equi}}$ of algorithms with **equidistant** time discretization:

- choose finite set $\mathcal{I} \subset \mathbb{N}$: only evaluate β_i with $i \in \mathcal{I}$.
- ▶ choose $n_i \in \mathbb{N}$: number of evaluations for β_i with $i \in \mathcal{I}$.
- ▶ choose nodes $t_{k,i} = \frac{k}{n_i} T$, $k = 1, \dots, n_i$, for β_i with $i \in \mathcal{I}$.
- ▶ choose $\phi : \mathbb{R}^N \rightarrow H$ measurable with $N = \sum_{i \in \mathcal{I}} n_i$.

Approximation:

$$\widehat{X}(T) = \phi(\beta_{i_1}(t_{1,i_1}), \dots, \beta_{i_1}(t_{n_{i_1},i_1}), \dots, \beta_{i_\ell}(t_{1,i_\ell}), \dots, \beta_{i_\ell}(t_{n_{i_\ell},i_\ell}))$$

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Clearly

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N -th minimal error

$$e_{\text{equi}}(N) = \inf \left\{ e(\widehat{X}(T)) ; \text{cost}(\widehat{X}(T)) \leq N \text{ and } \widehat{X}(T) \in \mathfrak{X}_{\text{equi}} \right\}.$$

Classes of Approximations and Minimal Errors

Class \mathfrak{X} of algorithms with **arbitrary** time discretization:

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- choose $n_i \in \mathbb{N}$: number of evaluations for β_i with $i \in \mathcal{I}$.
- ▶ choose nodes $0 < t_{1,i} < \dots < t_{n_i,i} \leq T$ for β_i with $i \in \mathcal{I}$.
- choose $\phi : \mathbb{R}^N \rightarrow H$ measurable with $N = \sum_{i \in \mathcal{I}} n_i$.

Approximation:

$$\widehat{X}(T) = \phi(\beta_{i_1}(t_{1,i_1}), \dots, \beta_{i_1}(t_{n_{i_1},i_1}), \dots, \beta_{i_\ell}(t_{1,i_\ell}), \dots, \beta_{i_\ell}(t_{n_{i_\ell},i_\ell}))$$

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for $\mathcal{I} = \{i_1, \dots, i_\ell\}$.

Clearly

$$\text{cost}(\widehat{X}(T)) = N.$$

N -th minimal error

$$e(N) = \inf \left\{ e(\widehat{X}(T)) ; \text{cost}(\widehat{X}(T)) \leq N \text{ and } \widehat{X}(T) \in \mathfrak{X} \right\}.$$

Classes of Approximations and Minimal Errors

N -th minimal errors

$$\begin{aligned}e_{\text{uni}}(N) &= \inf \left\{ e \left(\hat{X}(T) \right); \text{cost} \left(\hat{X}(T) \right) \leq N \text{ and } \hat{X}(T) \in \mathfrak{X}_{\text{uni}} \right\}, \\e_{\text{equi}}(N) &= \inf \left\{ e \left(\hat{X}(T) \right); \text{cost} \left(\hat{X}(T) \right) \leq N \text{ and } \hat{X}(T) \in \mathfrak{X}_{\text{equi}} \right\}, \\e(N) &= \inf \left\{ e \left(\hat{X}(T) \right); \text{cost} \left(\hat{X}(T) \right) \leq N \text{ and } \hat{X}(T) \in \mathfrak{X} \right\}.\end{aligned}$$

Clearly

$$e(N) \leq e_{\text{equi}}(N) \leq e_{\text{uni}}(N).$$

But are these minimal errors basically equal?

Results for Additive Noise

Theorem Henkel (2009)

Assume that $x_0 = 0$ and

$$\langle g(t) \cdot h_i, h_j \rangle^2 \asymp \frac{1}{|i-j|^\alpha + 1}$$

with $\alpha > 1$. Then

$$e_{\text{equi}}(N) \asymp e_{\text{uni}}(N) \asymp N^{-1/6}$$

and

$$N^{-1/2} \preceq e(N) \preceq \begin{cases} N^{-1/6} & \text{if } 1 < \alpha \leq \frac{5}{3}, \\ N^{-(\alpha-1)/4} & \text{if } \frac{5}{3} < \alpha < 2, \\ N^{-1/4} & \text{if } 2 \leq \alpha < \infty. \end{cases}$$

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Remark Limiting case $g = 1$, i.e. $\alpha \rightarrow \infty$:

$$e_{\text{uni}}(N) \asymp e_{\text{equi}}(N) \asymp N^{-1/6}$$

and

$$e(N) \asymp N^{-1/2},$$

see Müller-Gronbach, Ritter, Wagner (2007).

Results for Additive Noise

Remark

- ▶ Suboptimality of $\mathfrak{X}_{\text{equi}} \supset \mathfrak{X}_{\text{uni}}$ (at least), if $\alpha > \frac{5}{3}$.
- ▶ Optimal time discretizations are quantiles of the density $\exp(-\frac{\mu_j}{3}(T-t))$, i.e.

$$\int_0^{s_{k,j}} \exp\left(-\frac{\mu_j}{3}(T-t)\right) dt = \frac{k}{\nu_j} \int_0^T \exp\left(-\frac{\mu_j}{3}(T-t)\right) dt$$

for $j \in \mathcal{J} \subset \mathbb{N}$, $\mu_j = \pi^2 j^2$, $\nu_j \in \mathbb{N}$, $k = 1, \dots, \nu_j$ and $\{t_{1,i}, \dots, t_{n,i}\} = \bigcup_{j \in \mathcal{J}} \{s_{1,j}, \dots, s_{\nu_j,j}\}$ for every β_i .

Related Results

Equation

$$\begin{aligned}\frac{\partial}{\partial t}X(t, u) &= \frac{\partial^2}{\partial u^2}X(t, u) + a(t, u, X(t, u)) \\ &\quad + \sigma(t, u, X(t, u))\frac{\partial^2}{\partial t\partial u}W(t, u), \\ X(t, 0) &= X(t, 1) = 0, \quad t \in [0, T], \\ X(0, u) &= x_0(u), \quad u \in [0, 1],\end{aligned}$$

with regularity conditions for a , σ and x_0 .

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Theorem Gyöngy (1999)

With \widehat{X}_N , based on an equidistant N -point grid in $[0, T] \times [0, 1]$ to evaluate the Brownian sheet \mathcal{W} ,

$$\sup_{u \in [0, 1]} \left(\mathbb{E} |X(t, u) - \widehat{X}_N(t, u)|^2 \right)^{1/2} \preceq N^{-1/6}$$

for all $p \geq 1$ and $t \in (0, T]$.

Related Results

Equation

$$\begin{aligned}\frac{\partial}{\partial t}X(t, u) &= \frac{\partial^2}{\partial u^2}X(t, u) + X(t, u)\frac{\partial^2}{\partial t\partial u}\mathcal{W}(t, u), \\ X(t, 0) &= X(t, 1) = 0, \quad t \in [0, T], \\ X(0, u) &= 1, \quad u \in [0, 1].\end{aligned}$$

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Theorem *Davie, Gaines* (2001)

$$e(N) \asymp N^{-1/6}.$$

Related Results

Errorcriterion

$$\epsilon(\widehat{X}) = \left(\mathbb{E} \int_0^T \|X(t) - \widehat{X}(t)\|_H^2 dt \right)^{1/2},$$

i.e. approximation of $X(t)$ for every $t \in [0, T]$, see the following talk.

Summary and Outlook

Summary

At a fixed time point, for additive noise with decay condition

$$e_{\text{equi}}(N) \asymp e_{\text{uni}}(N) \asymp N^{-1/6}$$

and

$$N^{-1/2} \asymp e(N) \asymp \begin{cases} N^{-1/6} & \text{if } 1 < \alpha \leq \frac{5}{3}, \\ N^{-(\alpha-1)/4} & \text{if } \frac{5}{3} < \alpha < 2, \\ N^{-1/4} & \text{if } 2 \leq \alpha < \infty. \end{cases}$$

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Outlook

- ▶ Multiplicative noise.
- ▶ Sharp upper terms for $e(N)$.